# Using CDCAC for SMT inquiries with special constraints 

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## Introduction

- Recalling what SMT solvers do.
- Recalling CDCAC.
- Can we guide the search process in CDCAC?


## SMT solvers solve questions such as the following

## Let

- $x_{1}, \ldots, x_{n}$ be variables,
- $\mathbb{A}$ the set of real algebraic numbers $(\mathbb{Q} \subset \mathbb{A} \subset \mathbb{R})$,
- $f_{1}, \ldots, f_{m} \in \mathbb{A}\left[x_{1}, \ldots, x_{n}\right]$ some polynomials,
- $\prec_{1}, \ldots, \prec_{m}$ some relations from the set $\{=, \neq,<, \leq,>, \geq\}$,
- $\Phi\left(B_{1}, \ldots, B_{m}\right)$ a Boolean formula.

Then

$$
\stackrel{?}{\exists}\left(x_{1}, \ldots, x_{n}\right) \in \mathbb{R}^{n} \text { such that }\left.\Phi\left(B_{1}, \ldots, B_{m}\right)\right|_{B_{1}=f_{1} \prec_{1} 0, \ldots, B_{m}=f_{m} \prec_{m} 0} .
$$

## Example

$$
\stackrel{?}{\exists} x, y \in \mathbb{R} \text { such that } x>0 \wedge y>0 \wedge x^{2}+y^{2} \geq 1
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\left.B_{1} \wedge B_{2} \wedge B_{3}\right|_{B_{1}=(x>0), B_{2}=(y>0), B_{3}=\left(x^{2}+y^{2} \geq 1\right)}
$$

## Lazy SMT



## Example

Input formula: $\left(x^{2}+y^{2}<1 \wedge y>2\right) \vee\left(x^{2}+y^{2}<1 \wedge x>0\right)$
Lazy SMT sent two enquires to the theory solver, all in the form of conjunction of polynomial constraints.
1- $x^{2}+y^{2}<1 \wedge y>2 \wedge x \leq 0$ which is UNSAT.
2- $x^{2}+y^{2}<1 \wedge y \leq 2 \wedge x>0$ which is SAT with $(x, y)=\left(\frac{1}{2}, 0\right)$.

## CAD as the theory solver

## Example

$$
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## CDCAC as the theory solver

## CDCAC = Conflict Driven search using Cylindrical Algebraic Coverings

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$$

| $\begin{aligned} & c_{2}: 0<y \\ & c_{3}: 0 \leq y^{2}+3 \end{aligned}$ | $y$ $\left\{c_{2}\right\}$ | $\begin{array}{cc} (0, \infty) \\ \vdots & 2 \\ \vdots & \\ \vdots & \\ \vdots & 0\} \\ \vdots \end{array}$ |
| :---: | :---: | :---: |
|  | $\left\{c_{2}\right\}$ | $\begin{gathered} (-\infty, 0) \\ -2 \\ \Rightarrow y=2 \end{gathered}$ |

## CDCAC as the theory solver

## CDCAC $=$ Conflict Driven search using Cylindrical Algebraic Coverings

## Example

$$
x>0 \wedge y>0 \wedge x^{2}+y^{2} \geq 1 \quad \text { SAT },(x, y)=(2,2)
$$



## Comparing the theory solvers

## CDCAC vs CAD

```
[> CodeTools:-Usage( CDCAC ( [ x^2 + y^2 > 1, y > x^2 + 1 ], [ x, y ] ) );
memory used=4.32MiB, alloc change =32.00Mib, cpu time =63.00ms, real time =60.00ms, gc time=0ns
    $*
                                [true, [0, 3], []]
Full CAD.
restart:
    with (RegularChains:-SemiAlgebraicSetTools) :
    CodeTools:-Usage (CylindricalAlgebraicDecompose ([x^2+(y+2)^2-1,y-x^2-1],RegularChains:-PolynomialRing ([x,y]), output=cadcell)) :
memory used=6.58MiB, alloc change=32.00MiB, cpu time=78.00ms, real time=81.00ms, gc time=0ns
> restart:
    with(RegularChains:-SemiAlgebraicSetTools):
```



```
memory used=17.76MiB, alloc change=97.00Mib, cpu time=156.00ms, real time=278.00ms, gc time=0ns
> restart:
    with (RegularChains:-SemiAlgebraicSetTools):
    CodeTools:-Usage (CylindricalAlgebraicDecompose([x^2+(y+2)^2-1, y-x^2-1],RegularChains:-PolynomialRing ([x,y]), output=list)):
memory used=6.23Miв, alloc change = 32.00Mib, cpu time=63.00ms, real time=82.00ms,gc time=0ns
```


## Special request 1

## Guiding the CDCAC search

Did you notice the directions of movements in the CDCAC search steps? Another Example: $x^{2} y-5 x^{2}-5 x y-2 x-14 y-7>0, x^{2}+y^{2}-400<0$.


## Special request 1

## Guiding the CDCAC search

Did you notice the directions of movements in the CDCAC search steps? Another Example: $x^{2} y-5 x^{2}-5 x y-2 x-14 y-7>0, x^{2}+y^{2}-400<0$.


Current implementation returns $(0,-6)$.

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But what if we wanted a solution near a given point (the black point)?

## Special request 2

## Finding components of specific dimension

Did you notice the decompositions in each layer in CDCAC algorithm? There were closed (singleton sets) and open cells (open intervals). If the user wants a point from a solution component of dimension $d<n$, then we can avoid lifting up the partial points where there not enough closed cells used in its previous layers. For example we are at variable $x_{d+1}$ and all previous variables have picked up values from open intervals, so now we can ignore the open intervals in this layer.


## References

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Thank you for listening.
al-Khwarizmi

